

3008 Rec'd PCT/PTO 30 JAN 2001

FORM PTO-1390
REV. 5-93

US DEPARTMENT OF COMMERCE
PATENT AND TRADEMARK OFFICE

ATTORNEYS DOCKET NUMBER
P00,2004

**TRANSMITTAL LETTER TO THE UNITED STATES
DESIGNATED/ELECTED OFFICE (DO/EO/US)
CONCERNING A FILING UNDER 35 U.S.C. 371**

U.S. APPLICATION NO. (if known, see 37 CFR 1.5)

09/744829

INTERNATIONAL APPLICATION NO.
PCT/DE99/02016

INTERNATIONAL FILING DATE
01 JULY 1999

PRIORITY DATE CLAIMED
03 AUGUST 1998

TITLE OF INVENTION

METHOD FOR RE-ROUTING DATA PACKETS ONTO AN ALTERNATIVE NETWORK

APPLICANT(S) FOR DO/EO/US

CHRISTIAN PREHOFER

Applicant herewith submits to the United States Designated/Elected Office (DO/EO/US) the following items and other information:

1. ☒ This is a **FIRST** submission of items concerning a filing under 35 U.S.C. 371.
2. ☐ This is a **SECOND** or **SUBSEQUENT** submission of items concerning a filing under 35 U.S.C. 371.
3. ☒ This express request to begin national examination procedures (35 U.S.C. 371(f)) at any time rather than delay.
4. ☒ A proper Demand for International Preliminary Examination was made by the 19th month from the earliest claimed priority date.
5. ☒ A copy of International Application as filed (35 U.S.C. 371(c)(2)).
 - a. ☒ is transmitted herewith (required only if not transmitted by the International Bureau).
 - b. ☐ has been transmitted by the International Bureau.
 - c. ☐ is not required, as the application was filed in the United States Receiving Office (RO/US)
6. ☒ A translation of the International Application into English (35 U.S.C. 371(c)(2)).
7. ☒ Amendments to the claims of the International Application under PCT Article 19 (35 U.S.C. §371(c)(3))
 - a. ☐ are transmitted herewith (required only if not transmitted by the International Bureau).
 - b. ☐ have been transmitted by the International Bureau.
 - c. ☐ have not been made; however, the time limit for making such amendments has NOT expired.
 - d. ☒ have not been made and will not be made.
8. ☐ A translation of the amendments to the claims under PCT Article 19 (35 U.S.C. 371(c)(3)).
9. ☒ An oath or declaration of the inventor(s) (35 U.S.C. 371(c)(4)) - **unsigned**.
10. ☒ A translation of the annexes to the International Preliminary Examination Report under PCT Article 36 (35 U.S.C. 371(c)(5)).

Items 11. to 16. below concern other document(s) or information included:

11. ☒ An Information Disclosure Statement under 37 C.F.R. 1.97 and 1.98; (PTO 1449, Prior Art, Search Report, References).
12. ☐ An assignment document for recording. A separate cover sheet in compliance with 37 C.F.R. 3.28 and 3.31 is included.
(SEE ATTACHED ENVELOPE)
13. ☒ Amendment "A" Prior to Action.
 - ☐ A SECOND or SUBSEQUENT preliminary amendment.
14. ☒ A substitute specification and substitute specification mark-up.
15. ☐ A change of address letter attached to the Declaration.
16. ☒ Other items or information:
 - a. ☒ Submission of Drawings and Drawing Changes
 - b. ☒ EXPRESS MAIL #EL655302789US dated January 30, 2001

U.S. APPLICATION NO. (if known, see 37 C.F.R. 1.51) <div style="font-size: 24pt; font-weight: bold; margin-top: 5px;">09/744829</div>		INTERNATIONAL APPLICATION NO. PCT/DE99/02016		ATTORNEY'S DOCKET NUMBER P00,2004	
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17. <input checked="" type="checkbox"/> The following fees are submitted: <div style="margin-top: 10px;"> BASIC NATIONAL FEE (37 C.F.R. 1.492(a)(1)-(5): Search Report has been prepared by the EPO or JPO \$860.00 International preliminary examination fee paid to USPTO (37 C.F.R. 1.482) .. \$690.00 No international preliminary examination fee paid to USPTO (37 C.F.R. 1.482) but international search fee paid to USPTO (37 C.F.R. 1.445(a)(2)) \$710.00 Neither international preliminary examination fee (37 C.F.R. 1.482) nor international search fee (37 C.F.R. 1.445(a)(2)) paid to USPTO \$1000.00 International preliminary examination fee paid to USPTO (37 C.F.R. 1.482) and all claims satisfied provisions of PCT Article 33(2)-(4) \$ 100.00 <div style="text-align: right; margin-top: 10px;">ENTER APPROPRIATE BASIC FEE AMOUNT =</div> </div>				CALCULATIONS		PTO USE ONLY	
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Surcharge of \$130.00 for furnishing the oath or declaration later than <input type="checkbox"/> 20 <input type="checkbox"/> 30 months from the earliest claimed priority date (37 C.F.R. 1.492(e)).				\$			
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Claims	Number Filed	Number Extra	Rate			
Total Claims	12 - 20 =	0	X \$ 18.00	\$		
Independent Claims	01 - 3 =	0	X \$ 80.00	\$		
Multiple Dependent Claims			\$270.00 +	\$		
TOTAL OF ABOVE CALCULATIONS =				\$ 860.00		
Reduction by 1/2 for filing by small entity, if applicable. Verified Small Entity statement must also be filed. (Note 37 C.F.R. 1.9, 1.27, 1.28)				\$		
SUBTOTAL =				\$ 860.00		
Processing fee of \$130.00 for furnishing the English translation later than <input type="checkbox"/> 20 <input type="checkbox"/> 30 months from the earliest claimed priority date (37 CFR 1.492(f)).				+		
TOTAL NATIONAL FEE =				\$ 860.00		
Fee for recording the enclosed assignment (37 C.F.R. 1.21(h). The assignment must be accompanied by an appropriate cover sheet (37 C.F.R. 3.28, 3.31). \$40.00 per property				+		
TOTAL FEES ENCLOSED =				\$ 860.00		
				Amount to be refunded	\$	
				charged	\$	

a. ☒ A check in the amount of \$ 860.00 to cover the above fees is enclosed.

b. ☐ Please charge my Deposit Account No. _____ in the amount of \$ _____ to cover the above fees.
A duplicate copy of this sheet is enclosed.

c. ☒ The Commissioner is hereby authorized to charge any additional fees which may be required, or credit any
overpayment to Deposit Account No. 50-1519. A duplicate copy of this sheet is enclosed.

NOTE: Where an appropriate time limit under 37 C.F.R. 1.494 or 1.495 has not been met, a petition to revive (37 C.F.R. 1.137(a) or (b)) must be
filed and granted to restore the application to pending status.

SEND ALL CORRESPONDENCE TO:

SCHIFF HARDIN & WAITE
PATENT DEPARTMENT
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SIGNATURE

Steven H. Noll
NAME

28,982
Registration Number

09/744829

JC02 Rec'd PCT/PTO 30 JAN 2001

-1-

BOX PCT

IN THE UNITED STATES DESIGNATED/ELECTED OFFICE
IN THE UNITED STATES PATENT AND TRADEMARK OFFICE
UNDER THE PATENT COOPERATION TREATY - CHAPTER II

**AMENDMENT "A" PRIOR TO ACTION AND
SUBMISSION OF SUBSTITUTE SPECIFICATION**

APPLICANT: PREHOFER, Christian
ATTORNEY DOCKET NO: P00,2004
INTERNATIONAL APPLICATION NO: PCT/DE99/02016
INTERNATIONAL FILING DATE: 01 JUL 1999
INVENTION: METHOD FOR RE-ROUTING DATA
PACKETS TO ANOTHER NETWORK

Assistant Commissioner for Patents
Washington, DC 20231

Sir:

Applicant herewith submits an amendment and substitute specification in the above-reference PCT application, and respectfully requests entry of same prior to examination in the United States National Examination Stage.

IN THE SPECIFICATION

Cancel the specification as filed, and insert therefore the substitute specification provided herewith.

FILED - 09/744829

IN THE CLAIMS

Cancel claims 1 - 12 as filed, and insert therefore new claims 13 - 24 as follows:

- - What is Claimed is:

13. A method for re-routing data packets of a packet-switching network onto at least one alternate network capable of assuring a quality requested by a network user, the packet-switching network and the at least one alternate network form sub-networks of a network over which data packets can be transmitted, including at least one source node and at least one destination node that are each respectively one of either directly or indirectly connected to an access node via at least one intermediate node, the access node being capable of setting up a connection both to the packet-switching network and to one of the at least one alternate network, the method comprising:

identifying only by a respective bit pattern known to the access node the data packets to be routed via an alternate network in the source node by a bit pattern known to the access node that is connected to the source node either directly or indirectly via at least one intermediate node;

recognizing said known bit pattern upon arrival of such data packets in the access node; and

re-routing the data packets identified with only the known bit pattern onto an alternate network.

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14. A method according to claim 13, further comprising the step of using a filter in the access node to check data packets arriving from a source node for the known bit pattern; and

initiating the re-routing of the data packets identified with this bit pattern onto an alternate network when a known bit pattern is recognized.

15. A method according to claim 13, further comprising the step of connecting to the source node one of either directly or indirectly via at least one intermediate node containing a table for determining traffic paths into which the function of the filter is integrated, the table additionally contains bit patterns that can produce a re-routing of the data packet identified with such bit patterns onto an alternate network.

16. A method according to claim 15, further comprising the step of locating the known bit pattern in the header of a data packet to be routed via the alternate network.

17. A method according to claim 16, further comprising the step of using the same bit pattern in at least one source node regardless of the respectively requested quality.

18. A method according to claim 16, further comprising the step of using in at least one source node, bit patterns corresponding to the respectively requested quality.

19. A method according to claim 18, further comprising the step of using bit pattern of a data packet to produce a re-routing thereof onto at least one alternate network corresponding to the bit pattern with a specific quality.

20. A method according to claim 18, further comprising the step of using each recognized bit pattern of a data packet to produce a re-routing thereof onto at least one alternate network with a quality corresponding to the recognized bit pattern.

21. A method according to claim 20, further comprising the step of preventing the re-routing of the data packet onto at least one alternate network, if after recognition of such a bit pattern of a data packet to be routed via the at least one alternate network in such an access node, the at least one alternate network cannot offer the quality corresponding to the bit pattern.

22. A method according to claim 21, further comprising the steps of having the at least one source node send the data packets to communicate a message via the packet-switching network to at least one destination node with respect to the data packets to be routed via the at least one alternate network; and waiting for an acknowledge from the at least one destination node.

23. A method according to claim 22, further comprising the step of having the access node connected to the at least one source node send a message with

respect to the assured quality requested by the at least one source node to the network node the at least one alternate network; and

waiting for an acknowledge thereof.

24. A method according to claim 23, further comprising the step of reserving for data packets to be transmitted with an assured quality, at least one logical channel of the packet switching network, in a network constellation in which the at least one alternate network is formed. - -

IN THE ABSTRACT

Cancel the Abstract as filed and insert therefore on a separate page the following Abstract of the Disclosure:

- - ABSTRACT OF THE DISCLOSURE

Data packets of a packet-switching network for which transmission of a requested quality is to be assured, are routed via an alternate network. In a source node, such data packets are respectively identified by a bit pattern known to an access node, which is connected to the source node, either directly or indirectly, via at least one intermediate node. The bit pattern is recognized upon arrival in the access node, and as a result of identifying the data packets with the known bit pattern, the data packets are re-routed onto the alternate network. - -

REMARKS

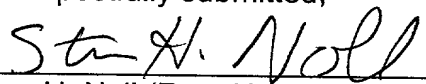
A substitute specification combining the specification and substitute pages, and an Abstract of the Disclosure are provided herewith, which make editorial changes in order to conform to standard US practice. A marked-up copy of the specification is also provided reflecting the changes made.

In addition, the claims as filed have been canceled and replaced by new claims that more clearly set forth the subject matter of Applicant's invention.

No new matter has been inserted into the application.

Applicant submits that this application is in proper condition for examination in the United States National Examination Stage, which action is respectfully requested.

Respectfully submitted,



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RECEIVED

Substitute Specification:

METHOD FOR RE-ROUTING DATA PACKETS ONTO AN ALTERNATIVE NETWORK

BACKGROUND OF THE INVENTION

Field of the Invention

The present invention is generally directed to telecommunication networks. In particular, the present invention is directed to methods for re-routing data packets on a data switching network.

Discussion of the Related Art

In a packet-switching network, such as for example the Internet, which is usually composed of a plurality of sub-networks, data packets from a source node are potentially transmitted to a destination node via a plurality of intermediate and/or access nodes of the individual sub-networks. In addition to containing information, the data packets particularly contain a destination address. The intermediate and/or access nodes contain what is referred to as a routing table for determining a traffic path in which each destination address of a data packet has a destination address of that node such as an intermediate node or an access node, to which the data packet is forwarded allocated to it. Thus, when a data packet arrives at such a node, the data packet is forwarded to the node corresponding to its destination address entry in the routing table.

The data transmission in such a packet-switching network is normally connectionless, i.e. the data packets with identical source and destination address

are transported independently of one another, so that neither the sequence nor a delivery of the data packets at the destination node is guaranteed, such as under OSI Layer 3 Protocol. However, a quality of the transmission of data packets between source and destination node such as certain bandwidth, delay times and a specific throughput, can not be promised.

Video transmission services, such as video on demand and various telephone services, such as voice over IP in the Internet, require a dependable and fast data transmission with an assured quality.

An "Internet Draft" document authored by K. Nichols and S. Blake that was published by the Internet Engineering Task Force in February 1998 (Internet site: <http://www.ietf.org/internet-drafts/draft-nichols-dsopdef-00.txt>) proposes a method that enables an accelerated transmission of data packets from a source node to a destination node.

For data packets to be forwarded especially fast, specific bits of what is referred to as the TOS (type of service) byte in the header part of such a data packet are set. Such data packets can be divided into various classes on the basis of the bits that are set in the TOS byte. According to their class, the data packets identified with the set bits are given privileged handling in the intermediate nodes via which such data packets are transmitted from the source node to the destination node, as a result an accelerated forwarding to the next node, intermediate or destination node is achieved.

A critical disadvantage of this method is that the privileged handling in the forwarding of the data packets identified with the set bits leads to considerable delays in the forwarding of the data packets to be transmitted normally.

Moreover, it is not only the destination addresses but also the respective TOS bytes that must be taken into consideration in every intermediate node when forwarding the data packets.

Another "Internet Draft" document authored by Y. Bernet, R. Yavatkar, P. Ford, F. Baker and L. Zhang that was published by the Internet Engineering Task Force in March 1998 (Internet site: <http://ietf.org/internet-drafts/draft-bernet-intdiff-00.txt>) presents a method that, with the assistance of the above-explained method, combines a plurality of demands made of the Internet for a guaranteed quality for the transmission of data packets into classes. A corresponding quality for the transmission of data packets is assured according to such a class.

Since a required quality for the transmission of data packets with a privileged handling thereof on the basis of the aforementioned, set bits in the TOS byte is assured, this method can be reduced to the initially explained approach. However, this method too suffers from the disadvantages mentioned above.

Another approach to offering the user of a video transmission service of a bandwidth suitable for the transmission of data packets is disclosed by US Patent 5,732,078.

That patent discloses an arrangement of an access node to the Internet that assures a bandwidth requested by the user for the transmission of data packets by re-routing data packets onto an alternate network. The re-routing of the data packets requires a user to request a specific bandwidth for the transmission of data packets from his user terminal device to a destination node. Then, the access node to which the user terminal device is connected sets up a point-to-point connection to

the access node to which the destination node is connected via the alternate network offering the requested bandwidth.

To re-route the data packets for whose transmission the user requests an assured bandwidth onto the alternate network, an existing routing table in the access node to which the user terminal device is connected is finally modified such that, in addition to containing the respective destination addresses of the nodes to which data packets are respectively forwarded, it also contains the source addresses of the data packets to be re-routed due to a quality requested by the user. According to a modification of the method presented in the 5,732,078 patent mentioned above, a connection-individual or a transmission-individual particular, also known as an application port number, is entered into the routing table.

On the basis of the additionally stored source address and the connection-individual or transmission-individual particular in the routing table, the data packets arriving at such an access node can be selected according to whether they are routed over the ordinary Internet or via the alternate network.

This method is very involved since a separate connection via the alternate network must be set up for every transmission of data packets with a requested bandwidth that is initiated by a user.

Additionally, the routing table must be modified in every access node to which user terminal devices are connected, namely after every transmission of data packets with a specific bandwidth initiated by a user.

Also to be considered is that all data packets arriving at such an access node are examined for their source address in order for the connection-individual or

transmission-individual particular, with reference to the source address, to be re-routing onto the alternate network.

In their Master Thesis, "Support Qos in IP over ATM", June 1997 (1997-06), National Taiwan University of Science and Technology, Taipei, Taiwan, (available online), Gung-Chou Lai and Ruay-Shiung Chang state that a determination regarding the path over which a data packet is to be routed in a switching system can be made in the switching system on the basis of a field TOS (type of service) in the header of data packets to be switched or transmitted. However, their presentation has nothing to do with re-routing of data packets of a packet-switching network onto at least one alternate network. On the contrary, the data switching or transmissions uses the same switching system.

SUMMARY OF THE INVENTION

therefore, it is an object of the present invention to provide a method for switching data packets of a specific quality in a packet switching network onto an alternate network.

It is another object of the present invention to provide a method for switching data packets of a specific quality in a packet switching network onto an alternate network, wherein a routing table remains unaffected by ongoing modifications for the re-routing of the data packets.

It is a further object of the present invention to provide a method for switching data packets of a specific quality in a packet switching network onto an alternate network, wherein data packets need only be examined for a known bit pattern in the access node connected to a source node.

It is an additional object of the invention to provide a method for switching data packets of a specific quality in a packet switching network onto an alternate network, wherein data packets arriving in an access node directly or indirectly connected to a source node proceeding from the source node are checked with a filter for a bit pattern known to the access node.

It is yet another object of the invention to provide a method for switching data packets of a specific quality in a packet switching network onto an alternate network, wherein the data packets are identified in their source node by a bit pattern known to the access node connected to the source node either directly or indirectly via at least one intermediate node.

BRIEF DESCRIPTION OF THE DRAWINGS

FIGURE 1 shows a data switching network according to the present invention.

DETAILED DESCRIPTION OF THE PREFERRED EMBODIMENTS

Accord to Figure 1, two source nodes U1, U2 of a plurality of conceivable source nodes and one of many possible destination nodes Z. The two source nodes U1, U2 are connected to an access node ZK1 and the destination node Z is connected to an access node ZK2, either directly or indirectly via respectively one or more intermediate nodes ZW1, ZW2. The intermediate node ZW1 is shown with broken lines between U1 and ZK1, and the intermediate node ZW2 is shown with broken lines between ZK2 and Z. The access nodes ZK1 and ZK2 belong to a packet-switching network PN and respectively form an access to one or more alternate networks AN1 through ANn at the same time. Such a packet-switching network could be the Internet for example, within which mainly variable-length data packets are transmitted. Or, such a packet switching network could also be an ATM

network (asynchronous transfer mode), wherein fixed-length data packets (ATM cells) are usually transported. Given the pre-condition that such an alternative network assures a requested quality for the data transmission, the alternative network can be formed by an arbitrary network type, for example a line-switching network, a packet-switching network or an ATM network. Such an alternate network can also be composed of at least one logical channel of the packet-switching network PN.

Normally, data packets are sent by an access node, for example ZK1, to the packet-switching network PN from the source node either directly (for example, proceeding from U2) or indirectly (for example, proceeding from U1), via one or more intermediate nodes, for example ZW1. From the access node, they are sent via the packet-switching network to another access node, for example ZK2, and are sent from the latter to a destination node Z either directly or via one or more intermediate nodes, for example ZW2. A transmission of data packets in the opposite direction is likewise conceivable, i.e. with the destination node as source node and with the source node as destination node.

In view of an overall telecommunication network, the indicated source node and the destination node can be respectively interpreted as intermediate node or end node of the telecommunication network. Such end nodes can thereby be viewed as a computer of a service vendor or as a data terminal device of a user.

In the present example, a user who would like to have his data terminal device, for example U1, receive an Internet service, for example video on demand, from a computer, for example Z, of a service vendor. For the transmission of data packets of such a service, a certain quality in the form of a specific bandwidth must

be assured at the network side. The bandwidth needed for the service can be offered by re-routing the data packets belonging to such a service onto an alternate network that assures the requested quality.

Accordingly, the source node, for example U1, sends a message to the access node ZK1 either directly or indirectly via at least one intermediate node, for example ZW1, this message containing a request in the form of re-routing data packets for whose transmission a specific quality, for example bandwidth, is required. In the form of a message, the access node communicates the bit pattern with which the data packets to be routed onto the alternate network are identified to the source node U1 and additionally acknowledges the message with the request sent from the source node U1. Further, the access node sends a message about the quality requested by the source node to the network node (not shown in Figure 1) of an alternate network, for example AN1, and potentially waits for the acknowledge thereof. The source node U1 identifies the data packets that are to be transmitted via the alternate network with an assured quality with the bit pattern obtained from the access node.

As an alternative thereto, such a bit pattern can be known to the source node U1 according to a corresponding implementation without having to be previously informed thereof by the access node. Such a bit pattern is normally located in the header of such a data packet. In conjunction with the Internet, the bit pattern can be found in what is referred to as the TOS byte.

Before the source node U1 sends the data packets, for whose transmission a requested quality is to be assured and which are therefore to be routed via an alternate network, in the direction of the access node, it is possible that the source

node communicates a message with respect to the data packets to be routed via the alternate network to its destination node via the ordinary packet-switching network PN and potentially waits for an acknowledge from the destination node.

In the access node, the data packets arriving from the source node U1 are checked for the bit pattern with a filter (not shown in the Figure 1). The data packets wherein the bit pattern was recognized are re-routed onto the alternate network. Alternatively, the function of the filter can be integrated into the routing table as being present in the access node, in that the bit patterns that produce a re-routing of a data packet identified with the bit pattern onto the alternate network, are additionally entered into the routing table.

Parallel to the above-described scenario with respect to the source node U1, the same scenario can be initiated by one or more other source nodes, for example U2.

Regardless of the requested quality, the source node U2 possibly employs the same bit pattern for identifying the data packets that are to be sent via an alternate network, for example ANn. Alternatively thereto, the source node U2 can use a bit pattern that corresponds to the requested quality for identifying the data packets to be routed via the alternate network.

In this way, the data packets sent in the direction of the access node from, a plurality of source nodes can be combined into classes according to the quality required for their transmission.

When the data packets to be re-routed in the access node are re-routed onto a plurality of alternate networks, one bit pattern is defined for each alternate

network. In this way, the data packets can be re-routed onto an alternate network corresponding to the bit pattern that assures the requested quality.

When only one alternate network is available for re-routing the data packets identified with a bit pattern, the quality to be assured by the alternate network is offered by the bit pattern of the identified data packets.

The access node can deny re-routing onto the alternate network for data packets that are identified with a bit pattern that does not correspond to the quality offered by the alternate network.

Mutually independent actions of the above-described scenarios can also run in an arbitrary sequence.

The method of the present invention causes data packets of a packet switching network for whose transmission between their source node and their destination a specific quality is requested onto at least one alternate network. This re-routing is made possible in that the data packets to be routed via the alternate network are identified in their source node by a bit pattern known to the access node connected to the source node either directly or indirectly via at least one intermediate node. Upon arrival of the data packets in such an access node, the known bit patterns are recognized only as a result of themselves. Thus, the data packets identified with the known bit patterns are re-routed onto an alternate network.

A critical advantage of the method of the present invention is that the table present in an access node connected to a source node for determining the traffic paths, also known as the routing table, remains unaffected by ongoing modifications

for the re-routing of the data packets for whose transmission a specific quality has been requested.

Also, when data packets need only be examined for the known bit pattern in the access node connected to a source node.

The method of the present invention is not negatively affected by the traffic flow of the ordinary packet-switching network is not negatively affected by the re-routing of the data packets. By selecting a known bit pattern, a desired transmission quality can be assured.

Data packets arriving in an access node directly or indirectly connected to a source node proceeding from the source node are checked with a filter for a bit pattern known to the access node. When a known bit pattern has been recognized, the re-routing of the data packets identified with such a bit pattern onto an alternate network is initiated.

The method of the present invention is particularly beneficial particularly because all data packets arriving at the access node need not be examined for the known bit pattern, but only a significantly lower plurality of data packets, namely those coming from a source node. In addition, data packets that are sent from the destination node back to the source node, for example for the purpose of an acknowledgment, are re-routed onto the alternate network. Thus, the data packets are unintentionally sent in a circle.

The function of the filter is integrated in a table for determining the traffic paths, also known as the routing table, that is present in such an access node. This occurs with an additional entry of the known bit pattern, thereby producing a re-routing of a data packet identified with such a bit pattern onto an alternate network.

Implementation of the filter is simplified because the memory structures already present can be utilized to integrate the function of the filter into the routing table.

The known bit pattern is located in the header part of a data packet to be routed via an alternate network. As a result thereof, data packets can be designationally and examined for the known bit pattern.

The method of the present invention uses the same bit pattern for identifying data packets to be routed via an alternate network regardless of the requested quality. Thus, examination of such data packets for a known bit pattern is considerably simplified. Further, data packets for whose transmission a different transmission quality is requested are combined into a class, so that the alternate network need offer only one of the requested qualities for the transmission of these data packets.

As an alternative, the source nodes can also employ different bit patterns for identifying such data packets. These different bit patterns correspond to the respectively requested quality. This has the advantage that such data packets can be divided into classes according to their requested quality. The alternate network can thus offer the quality corresponding to a class for the transmission of the data packets.

Another embodiment of the method of the present invention provides that each recognized bit pattern of a data packet produces a re-routing thereof onto an alternate network with a requested quality that corresponds to the bit pattern. In other words, data packets of a class are re-routed onto one of the possible alternate networks that corresponds to the bit pattern of the class and offers the quality requested by the class.

In another embodiment of the present invention, each recognized bit pattern of a data packet can produce a re-routing thereof onto a single alternate network that assures the quality that corresponds to the bit pattern and, thus, to the class of the data packets.

The above-presented developments for dividing the data packets into classes corresponding to the requested quality offers the advantage that an alternate network need not explicitly offer the respectively requested quality for each transmission of data packets to be undertaken.

After recognition of such a bit pattern of a data packet to be routed via an alternate network in such an access node, the re-routing of the data packet onto an alternate network can be prevented when the alternate network cannot offer the quality corresponding to the bit pattern. In this way, such an access node can defend such an alternate network against data packets that lead to an overload.

In addition, the source node that intends to send data packets, communicates a message to its destination node via the packet-switching network with respect to the data packets to be routed via an alternate network and potentially waits for an acknowledge from the destination node. A secured connection setup of a point-to-point connection between the source and the destination node is achieved as a result thereof.

Further, the access node directly or indirectly connected to such a source node sends a message with respect to the quality to be assured as requested by the source node to the network node of an alternate network and potentially waits for an acknowledge therefrom. This procedure serves the purpose of dependably offering of the requested quality by the network node of the alternate network.

The method of the present invention can be applied to a network constellation wherein such an alternate network is formed, wherein that at least one logical channel of the packet-switching network is reserved for the data packets to be transmitted with an assured quality. For instance, an alternate network is composed of one or more logical channels of the packet-switching network. This represents an especially cost-beneficial solution since the need for additional connecting lines for the alternate network is eliminated.

Although modifications and changes may be suggested by those skilled in the to which this invention pertains, it is the intention of the inventor to embody with the patent warranted hereon all changes and modifications that may reasonably and properly come under the scope of their contribution to the art.

~~{Siemens AG}~~

New PCT application

26965-0762 (P-00,2004)

1998P02235WOUS

Inventor: Prehofer

Translation / January 11, 2001 / 1696(911) / 4300 words} **[Substitute
Specification:]**

METHOD FOR RE-ROUTING DATA PACKETS

~~{OF A PACKET SWITCHING NETWORK ONTO AT LEAST ONE}~~ **[ONTO AN]
ALTERNATIVE NETWORK**

[BACKGROUND OF THE INVENTION]

Field of the Invention

The present invention is generally directed to telecommunication networks. In particular, the present] ~~{The}~~ invention is directed to ~~{a method}~~ **[methods]** for re-routing data packets ~~{of}~~ **[on]** a data switching network ~~{onto at least one alternative network that assures a quality required for these data packets, according to the preamble of patent claim 1.}~~.

Discussion of the Related Art]

In a packet-switching network[,] such as~~{,}~~ for example~~{,}~~ the Internet, which is usually composed of a plurality of sub-networks, data packets from a source node

are potentially transmitted to a destination node via a plurality of intermediate and/or access nodes of the individual sub-networks. In addition to containing information, the data packets particularly contain a destination address. The intermediate ~~and/~~ or ~~{, respectively,}~~ access nodes contain what is referred to as a routing table for determining a traffic path in which each destination address of a data packet has a destination address of that node ~~{(i.e.)}~~ **[such as an] intermediate [node]** or ~~{, respectively,}~~ **[an] access node** ~~{,}~~ to which the data packet is forwarded allocated to it. ~~{When, thus,}~~ **[Thus, when]** a data packet arrives at such a node, the data packet is forwarded to the node corresponding to its destination address entry in the routing table.

The data transmission in such a packet-switching network is normally connectionless, i.e. the data packets with identical source and destination address are transported independently of one another, so that neither the sequence nor a delivery of the data packets at the destination node is guaranteed ~~{(OSI layer 3 protocol). A}~~ **[, such as under OSI Layer 3 Protocol. However, a]** quality of the transmission of data packets between source and destination node such as ~~{, for example, a}~~ certain bandwidth, delay times and a specific throughput **[,]** can ~~{therefore also}~~ not be promised.

~~{In particular, video}~~ **[Video]** transmission services ~~{(for example,)}~~ **[, such as]** video on demand ~~{}~~ and various telephone services ~~{(for example,)}~~ **[, such as]** voice over IP ~~{}~~ in the Internet **[,]** require a dependable and fast data transmission with an assured quality.

An "Internet Draft" document authored by K. Nichols and S. Blake that was published by the Internet Engineering Task Force in February 1998 (Internet site:

http://www.ietf.org/internet-drafts/draft-nichols-dsopdef-00.txt) proposes a method that enables an accelerated transmission of data packets from a source node to a destination node.

For data packets to be forwarded especially fast, specific bits of what is referred to as the TOS **[(type of service)]** byte in the header part of such a data packet are set. Such data packets can be divided into various classes on the basis of the bits that are set in the TOS byte. According to their class, the data packets identified with the set bits are given privileged handling in the intermediate nodes via which such data packets are transmitted from the source node to the destination node, as a result ~~{whereof, in particular,}~~ an accelerated forwarding to the next node~~{([,] intermediate or destination node{)}~~ is achieved.

A critical disadvantage of this method is ~~{to be seen therein}~~ that the privileged handling in the forwarding of the data packets identified with the set bits leads to considerable delays in the forwarding of the data packets to be transmitted ~~{}~~ normally~~{}~~.

Moreover, it is not only the destination addresses but also the respective TOS bytes that must be taken into consideration in every intermediate node when forwarding the data packets.

Another "Internet Draft" document authored by Y. Bernet, R. Yavatkar, P. Ford, F. Baker and L. Zhang that was published by the Internet Engineering Task Force in March 1998 (Internet site: <http://ietf.org/internet-drafts/draft-bernet-intdiff-00.txt>) presents a method that, with the assistance of the above-explained method, combines a plurality of demands made of the Internet for a guaranteed quality for the transmission of data packets

into classes. A corresponding quality for the transmission of data packets is assured according to such a class.

Since a required quality for the transmission of data packets with a privileged handling thereof on the basis of the aforementioned, set bits in the TOS byte is assured, this method can be reduced to the initially explained approach. ~~{The}~~ **[However, this method too suffers from the]** disadvantages ~~{that were already initially}~~ mentioned ~~{therefore remain}~~ **[above]**.

Another approach to offering the user of~~{, for example,}~~ a video transmission service **[of]** a bandwidth ~~{requested by the user}~~ **[suitable]** for the transmission of data packets is disclosed by US Patent ~~{Application}~~ 5,732,078.

~~{This}~~ **[That patent]** discloses an arrangement of an access node to the Internet that assures a bandwidth requested by the user for the transmission of data packets by re-routing data packets onto an alternate network. The re-routing of the data packets **[requires a user to request]** ~~{is thereby undertaken as follows: When a user requests}~~ a specific bandwidth for the transmission of data packets from his user terminal device to a destination node~~[. Then]~~, the access node to which the user terminal device is connected sets up a point-to-point connection to the access node to which the destination node is connected via the alternate network offering the requested bandwidth.

~~{For}~~ **[To]** re-~~{routing}~~ **[route]** the data packets for whose transmission the user requests an assured bandwidth onto the alternate network, an existing routing table in the access node to which the user terminal device is connected is finally modified such that, in addition to containing the respective destination addresses of the nodes to which data packets are respectively forwarded, it also contains the

source addresses of the data packets to be re-routed due to a quality requested by the user. [According to a modification of the method presented in the 5,732,078 patent mentioned above, a connection-individual or a transmission-individual particular, also know as an application port number, is entered into the routing table.]

On the basis of the additionally stored source address [and the connection-individual or transmission-individual particular] in the routing table, the data packets arriving at such an access node can be selected according to whether they are routed over the ordinary Internet or via the alternate network.

This method is very involved since a separate connection via the alternate network must be set up for every transmission of data packets with a requested bandwidth that is initiated by a user.

Additionally, the routing table must be modified in every access node to which user terminal devices are connected, namely after every transmission of data packets with a specific bandwidth initiated by a user.

Also to be considered ~~as further outlay is that all data packets incoming at such an access node are investigated for their source address, with reference whereto a re-routing onto the alternate network can be initiated.~~

~~The object of the invention is therefore comprised in developing a method of the species indicated in the preamble of patent claim 1 to the effect that it can be implemented with optimally little outlay and with the slightest possible effect on its environment.~~

~~The object is achieved by the features recited in the characterizing part of claim 1.~~

According to an advantageous development of the invention, the data packets arriving in an access node directly or indirectly connected to a source node proceeding from the source node are checked with a filter for a bit pattern known to the access node. When a known bit pattern has been recognized, then the re-routing of the data packets identified with such a bit pattern onto an alternate network is initiated. This procedure is beneficial particularly because it is no longer all data packets arriving at the access node but only a significantly lower plurality of data packets, namely those coming from a source node, that must be investigated for the known bit pattern. What is also thereby avoided is that data packets that are sent from the destination node back to the source node, for example for the purpose of an acknowledge, are re-routed onto the alternate network, as a result whereof the data packets are unintentionally sent in a circle.

Another advantageous development of the invention provides that the function of the filter is integrated in a table for determining the traffic paths (routing table) that is present in such an access node. This occurs with an additional entry of the known bit pattern that can produce a re-routing of a data packet identified with such a bit pattern onto an alternate network. In that memory structures that are already present can be utilized in the integration of the function of the filter into the routing table, the implementation outlay of such a filter is substantially reduced.

According to a useful development of the invention, such a known bit pattern is located in the header part of a data packet to be routed via an alternate network. As a result thereof, the data packets can be designationally and, thus, quickly investigated for the known bit pattern.

According to one development, the inventive method uses the respectively same bit pattern for identifying data packets to be routed via an alternate network, regardless of the respectively requested quality. The investigation of such data packets for a known bit pattern is considerably simplified as a result thereof. Further, the data packets for whose transmission a respectively different transmission quality is requested are combined into a class, so that the alternate network need offer only one of the requested qualities for the transmission of these data packets.

As an alternative to the aforementioned development, the source nodes can also employ different bit patterns for identifying such data packets, these different bit patterns corresponding to the respectively requested quality. This has the advantage that such data packets can be divided into classes according to their requested quality. The alternate network can thus offer the quality corresponding to a class for the transmission of the data packets.

In conjunction with the above-described development, another embodiment of the invention provides that each recognized bit pattern of a data packet produces a re-routing thereof onto an alternate network with a requested quality that corresponds to the bit pattern. In other words, data packets of a class are re-routed onto one of the possible alternate networks that corresponds to the bit pattern of the class and offers the quality requested by the class.

As an alternative to the aforementioned embodiment of the invention, each recognized bit pattern of a data packet can produce a re-routing thereof onto a single alternate network that assures the quality that corresponds to the bit pattern and, thus, to the class of the data packets.

~~network is formed in that at least one logical channel of the packet-switching network is reserved for the data packets to be transmitted with an assured quality. I.e., an alternate network is composed of one or more logical channels of the packet-switching network. This represents an especially cost-beneficial solution since additional connecting lines for the alternate network are thereby eliminated. An exemplary embodiment of the invention is explained in greater detail below with reference to a drawing.~~

~~The Figure shows a portion of a communication network on which the inventive method can be applied by way of example.~~

~~Accordingly, the Figure shows} [is that all data packets arriving at such an access node are examined for their source address in order for the connection-individual or transmission-individual particular, with reference to the source address, to be re-routing onto the alternate network.~~

In their Master Thesis, "Support Qos in IP over ATM", June 1997 (1997-06), National Taiwan University of Science and Technology, Taipei, Taiwan, (available online), Gung-Chou Lai and Ruay-Shiung Chang state that a determination regarding the path over which a data packet is to be routed in a switching system can be made in the switching system on the basis of a field TOS (type of service) in the header of data packets to be switched or transmitted. However, their presentation has nothing to do with re-routing of data packets of a packet-switching network onto at least one alternate network. On the contrary, the data switching or transmissions uses the same switching system.

SUMMARY OF THE INVENTION

therefore, it is an object of the present invention to provide a method for switching data packets of a specific quality in a packet switching network onto an alternate network.

It is another object of the present invention to provide a method for switching data packets of a specific quality in a packet switching network onto an alternate network, wherein a routing table remains unaffected by ongoing modifications for the re-routing of the data packets.

It is a further object of the present invention to provide a method for switching data packets of a specific quality in a packet switching network onto an alternate network, wherein data packets need only be examined for a known bit pattern in the access node connected to a source node.

It is an additional object of the invention to provide a method for switching data packets of a specific quality in a packet switching network onto an alternate network, wherein data packets arriving in an access node directly or indirectly connected to a source node proceeding from the source node are checked with a filter for a bit pattern known to the access node.

It is yet another object of the invention to provide a method for switching data packets of a specific quality in a packet switching network onto an alternate network, wherein the data packets are identified in their source node by a bit pattern known to the access node connected to the source node either directly or indirectly via at least one intermediate node.

BRIEF DESCRIPTION OF THE DRAWINGS

FIGURE 1 shows a data switching network according to the present invention.

DETAILED DESCRIPTION OF THE PREFERRED EMBODIMENTS

Accord to Figure 1,] two source nodes U1, U2 of a plurality of conceivable source nodes and one of many possible destination nodes Z. The two source nodes U1, U2 are connected to an access node ZK1 and the destination node Z is connected to an access node ZK2, either directly or indirectly via respectively one or more intermediate nodes ZW1, ZW2{, as indicated in the Figure by the} [. The intermediate node ZW1 [is] shown with broken lines between U1 and ZK1[, and the intermediate node ZW2 [is] shown with broken lines between ZK2 and Z. The access nodes ZK1 and ZK2 belong to a packet-switching network PN and respectively form an access to one or more alternate networks AN1 through ANn at the same time. Such a packet-switching network could{,} [be the Internet] for example, {be the Internet,} within which mainly variable-length data packets are transmitted{, or}[. Or, such a packet switching network] could also be an ATM network (asynchronous transfer mode), {within which} [wherein] fixed-length data packets (ATM cells) are usually transported. Given the pre-condition that {it} [such an alternative network] assures a requested quality for the data transmission, {such an alternate} [the alternative] network can be formed by an arbitrary network type, for example a line-switching network, a packet-switching network or an ATM network. Such an alternate network can also be composed of at least one logical channel of the packet-switching network PN.

Normally, data packets are sent {to} [by] an access node, for example ZK1, to the packet-switching network PN from the source node either directly (for example, proceeding from U2) or indirectly (for example, proceeding from U1)[,] via one or more intermediate nodes, for example ZW1. From the access node, they are

sent via the packet-switching network to another access node, for example ZK2, and are sent from the latter to a destination node Z either directly or via one or more intermediate nodes, for example ZW2. A transmission of data packets in the opposite direction is likewise conceivable, i.e. with the destination node as source node and with the source node as destination node.

In view of ~~the~~ **[an]** overall ~~communication~~ **[telecommunication]** network, the indicated source node and the destination node can be respectively interpreted as intermediate node ~~of~~ **[or]** end node of the ~~communication~~ **[telecommunication]** network. Such end nodes can thereby be viewed as **[a]** computer of a service vendor or as a data terminal device of a user.

In the present example, a user **[who]** would like to have his data terminal device, for example U1, receive an Internet service, for example video on demand, from a computer, for example Z, of a service vendor. For the transmission of data packets of such a service, a certain quality in the form of a specific bandwidth must be assured at the network side. The bandwidth needed for the service can be offered by re-routing the data packets belonging to such a service onto an alternate network that assures the requested quality.

Accordingly, the source node, for example U1, sends a message to the access node ZK1 either directly or indirectly via at least one intermediate node, for example ZW1, this message containing a request in the form **[of]** re-routing data packets for whose transmission a specific quality, for example bandwidth, is required. In the form of a message, the access node communicates the bit pattern with which the data packets to be routed onto the alternate network are identified to the source node U1 and additionally acknowledges the message with the request

sent from the source node U1. Further, the access node sends a message about the quality requested by the source node to the network node (not shown in ~~the Figure~~) **[Figure 1]** of an alternate network, for example AN1, and potentially waits for the acknowledge thereof. The source node U1 identifies the data packets that are to be transmitted via the alternate network with an assured quality with the bit pattern obtained from the access node.

As an alternative thereto, such a bit pattern can be known to the source node U1 according to a corresponding implementation without having to be previously informed thereof by the access node. Such a bit pattern is normally located in the header of such a data packet. In conjunction with the Internet, the bit pattern can be found in what is referred to as the TOS byte.

Before the source node U1 ~~now~~ sends the data packets~~{...}~~**[...]** for whose transmission a requested quality is to be assured and which are therefore to be routed via an alternate network ~~{...}~~**[...]** in the direction of the access node, it is possible that the source node communicates a message with respect to the ~~{...}~~**[...]** **[data packets]** to be routed via the alternate network to its destination node ~~{data packets}~~**[sic]** via the ordinary packet-switching network PN and potentially waits for an acknowledge from the destination node.

In the access node, the data packets arriving from the source node U1 are checked for the bit pattern with a filter (not shown in the Figure ~~{...}~~ **[1]**). The data packets wherein the bit pattern was recognized are re-routed onto the alternate network. Alternatively ~~(thereto)~~, the function of the filter can be integrated into the routing table ~~(to be conceived of)~~ as being present in the access node~~[...]~~ in that the

bit patterns that produce a re-routing of a data packet identified with the bit pattern onto the alternate network[,] are additionally entered into the routing table.

Parallel to the above-described scenario with respect to the source node U1, the same scenario can be initiated by one or more other source nodes, for example U2.

Regardless of the ~~{respectively}~~ requested quality, the source node U2 possibly employs the same bit pattern for identifying the data packets that are to be sent via an alternate network, for example ANn. Alternatively thereto, the source node U2 can use a bit pattern that corresponds to the requested quality for identifying the data packets to be routed via the alternate network.

In this way, the data packets sent in the direction of the access node from, ~~{under certain circumstances,}~~ a plurality of source nodes can be combined into classes according to the quality required for their transmission.

When the data packets to be re-routed in the access node ~~{can be}~~ **[are]** re-routed onto a plurality of alternate networks, one bit pattern is defined ~~{[...]}~~ **[for]** each alternate network. In this way, the data packets can be re-routed onto an alternate network corresponding to the bit pattern that assures the requested quality.

When only one alternate network is available for ~~{the}~~ re-routing ~~{of}~~ the data packets identified with a bit pattern, the quality to be assured by the alternate network ~~{can be}~~ **[is]** offered ~~{according to}~~ **[by]** the bit pattern of the identified data packets.

The access node can deny re-routing onto the alternate network for data packets that are identified with a bit pattern that does not correspond to the quality offered by the alternate network.

Mutually independent actions of the above-described scenarios can also run in an arbitrary sequence.

~~{Patent Claims~~

~~1. Method for re-routing} [The method of the present invention causes] data packets of a packet {switching network (PN)} [switching network for whose transmission between their source node and their destination a specific quality is requested] onto at least one alternate network{(AN1, ..., ANn) that assures a quality requested for these data packets, whereby the packet-switching network and the at least one alternate network form sub-networks of a communication network that is composed of at least one source node (U1, U2) and at least one destination node (Z) that are respectively connected to an access node (ZK1, ZK2)}[. This re-routing is made possible in that the data packets to be routed via the alternate network are identified in their source node by a bit pattern known to the access node connected to the source node]either directly or indirectly via at least one intermediate node{, said access node being capable of setting up a connection both to the packet-switching network as well as to}[. Upon arrival of the data packets in such an access node, the known bit patterns are recognized only as a result of themselves. Thus, the data packets identified with the known bit patterns are re-routed onto] an alternate network{, and between which (U1, U2; Z) data packets can be transmitted, characterized in that the data packets to be routed via an alternate network are identified in their source node by a bit pattern known to the access node (ZK1) that is connected to the source node either directly or indirectly via at least one intermediate node (ZW1);~~

~~said bit pattern being respectively recognized upon arrival of such data packets in the access node, as a result whereof a)~~].

A critical advantage of the method of the present invention is that the table present in an access node connected to a source node for determining the traffic paths, also known as the routing table, remains unaffected by ongoing modifications for the]re-routing of the data packets {identified with} [for whose transmission a specific quality has been requested.

Also, when data packets need only be examined for] the known bit pattern {onto an alternate network is initiated.} [in the access node connected to a source node.]

~~{2. Method according to claim 1, characterized in that a filter in such an access node checks the data packets arriving from a source node for a bit pattern known to the access node and, when a known bit pattern was recognized, initiates}~~
[The method of the present invention is not negatively affected by the traffic flow of the ordinary packet-switching network is not negatively affected by] the re-routing of the data packets {identified with this}]. By selecting a known bit pattern, a desired transmission quality can be assured.

Data packets arriving in an access node directly or indirectly connected to a source node proceeding from the source node are checked with a filter for a bit pattern known to the access node. When a known bit pattern has been recognized, the re-routing of the data packets identified with such a] bit pattern onto an alternate network [is initiated.

The method of the present invention is particularly beneficial particularly because all data packets arriving at the access node need not be

examined for the known bit pattern, but only a significantly lower plurality of data packets, namely those coming from a source node. In addition, data packets that are sent from the destination node back to the source node, for example for the purpose of an acknowledgment, are re-routed onto the alternate network. Thus, the data packets are unintentionally sent in a circle.

The] {

3. Method according to claim 1, characterized in that the access node connected to such a source node either directly or indirectly via at least one intermediate node contains a table for determining the traffic paths into which the } function of the filter is integrated {, in that is additionally contains the bit patterns that can produce} [in a table for determining the traffic paths, also known as the routing table, that is present in such an access node. This occurs with an additional entry of the known bit pattern, thereby producing] a re-routing of a data packet identified with such a bit pattern onto an alternate network. [Implementation of the filter is simplified because the memory structures already present can be utilized to integrate the function of the filter into the routing table.

The known] {4. Method according to one of the preceding claims, characterized in that such a} bit pattern is located in the header [part] of a data packet to be routed via an alternate network. [As a result thereof, data packets can be designationally and examined for the known bit pattern.

The method of the present invention uses] {5. Method according to one of the preceding claims, characterized in that} the same bit pattern {is employed by all source nodes regardless of the respectively requested quality.} [for identifying data packets to be routed via an alternate network regardless of the requested

quality. Thus, examination of such data packets for a known bit pattern is considerably simplified. Further, data packets for whose transmission a different transmission quality is requested are combined into a class, so that the alternate network need offer only one of the requested qualities for the transmission of these data packets.]

~~{6. Method according to one of the claims 1 through 4, characterized in that the source nodes employ bit patterns corresponding}~~ [As an alternative, the source nodes can also employ different bit patterns for identifying such data packets. These different bit patterns correspond] to the respectively requested quality. [This has the advantage that such data packets can be divided into classes according to their requested quality. The alternate network can thus offer the quality corresponding to a class for the transmission of the data packets.

Another embodiment of the method of the present invention provides] ~~{7. Method according to claim 6, characterized in}~~ that each recognized bit pattern of a data packet produces a re-routing thereof onto an alternate network [with a requested quality that corresponds] ~~{corresponding to the bit pattern with a specific quality.~~

~~8. Method according to claim 6, characterized in that each recognized bit pattern of a data packet produces a re-routing thereof onto an alternate network with a quality corresponding}~~ to the bit pattern. [In other words, data packets of a class are re-routed onto one of the possible alternate networks that corresponds to the bit pattern of the class and offers the quality requested by the class.

In another embodiment of the present invention, each recognized bit pattern of a data packet can produce a re-routing thereof onto a single alternate network that assures the quality that corresponds to the bit pattern and, thus, to the class of the data packets.

The above-presented developments for dividing the data packets into classes corresponding to the requested quality offers the advantage that an alternate network need not explicitly offer the respectively requested quality for each transmission of data packets to be undertaken.

After ~~{9. Method according to one of the preceding claims, characterized in that, after}~~ recognition of such a bit pattern of a data packet to be routed via an alternate network in such an access node, the re-routing of the data packet onto an alternate network can be prevented when the alternate network cannot offer the quality corresponding to the bit pattern. **[In this way, such an access node can defend such an alternate network against data packets that lead to an overload.**

In addition, ~~{10. Method according to one of the preceding claims, characterized in that}~~ the source node that intends to send data packets~~[,]~~ communicates a message **[to its destination node]** via the packet-switching network ~~{to its destination node}~~ with respect to the data packets to be routed via an alternate network and potentially ~~{wait}~~ **[waits]** for an acknowledge from the destination node. **[A secured connection setup of a point-to-point connection between the source and the destination node is achieved as a result thereof.**

Further, the access node directly or indirectly ~~{11. Method according to one of the preceding claims, characterized in that the access node}~~ connected to

such a source node ~~{either directly or indirectly via at least one intermediate node}~~
sends a message with respect to the ~~{assured}~~ quality **[to be assured as]** requested
by the source node to the network node of ~~{such}~~ an alternate network and
potentially waits for an acknowledge ~~{thereof.}~~ **[therefrom. This procedure serves
the purpose of dependably offering of the requested quality by the network
node of the alternate network.]**

~~{12. Method according to one of the preceding claims, characterized in that it
is}~~ **[The method of the present invention can be]** applied to a network
constellation wherein such an alternate network is formed ~~{in}~~ **[, wherein]** that at
least one logical channel of the packet-switching network is reserved for the data
packets to be transmitted with an assured quality. **[For instance, an alternate
network is composed of one or more logical channels of the]** ~~{Abstract
METHOD FOR RE-ROUTING DATA PACKETS OF A PACKET SWITCHING
NETWORK ONTO AT LEAST ONE ALTERNATIVE NETWORK
Data packets of a}~~ packet-switching network ~~{(PN) for whose transmission a
requested quality is to be assured are routed via at least one alternate network (AN
1, ..., ANn). In their source node (U1, U2), such data packets are respectively
identified by a bit pattern known [...]~~ access node ~~(ZK1) connected to the source
node either directly or indirectly via at least one intermediate node (ZW1), said bit
pattern being respectively recognized upon arrival of such data packets in the
access node, as a result whereof a re-routing of the data packets identified with the
known bit pattern onto an}~~ **[. This represents an especially cost-beneficial
solution since the need for additional connecting lines for the]** alternate
network is ~~{initiated.}~~ **[eliminated.]**

{Figure} [Although modifications and changes may be suggested by those skilled in the to which this invention pertains, it is the intention of the inventor to embody with the patent warranted hereon all changes and modifications that may reasonably and properly come under the scope of their contribution to the art.]

09/744829

JC02 Rec'd PCT/PTO 3 0 JAN 2001

Siemens AG

New PCT application

26965-0762 (P-00,2004)

1998P02235WOUS

Inventor: Prehofer

Translation / January 11, 2001 / 1696(911) / 4300 words

1/PRB

09/744829

JC02 Rec'd PCT/PTO 30 JAN 2001

- 1 -

**METHOD FOR RE-ROUTING DATA PACKETS OF A PACKET
SWITCHING NETWORK ONTO AT LEAST ONE ALTERNATIVE
NETWORK**

5 The invention is directed to a method for re-routing data packets of a data switching network onto at least one alternative network that assures a quality required for these data packets, according to the preamble of patent claim 1.

10 In a packet-switching network such as, for example, the Internet, which is usually composed of a plurality of sub-networks, data packets from a source node are potentially transmitted to a destination node via a plurality of intermediate and/or access nodes of the individual sub-networks. In addition to containing information, the data packets particularly contain a destination address. The intermediate or, respectively, access nodes contain what is referred to as a routing table for determining a traffic path in which each destination address of a data packet has a destination address of that node (i.e. intermediate or, respectively, access node) to which the data packet is forwarded allocated to it. When, thus, a data packet arrives at such a node, the data packet is forwarded to the node corresponding to its destination address entry in the routing table.

15 The data transmission in such a packet-switching network is normally connectionless, i.e. the data packets with identical source and destination address are transported independently of one another, so that neither the sequence nor a delivery of the data packets at the destination node is guaranteed (OSI layer 3 protocol). A quality of the transmission of data packets between source and destination node such as, for example, a certain bandwidth, delay times and a specific throughput can therefore also not be promised.

25 In particular, video transmission services (for example, video on demand) and various telephone services (for example, voice over IP) in the Internet require a dependable and fast data transmission with an assured quality.

An "Internet Draft" document authored by K. Nichols and S. Blake that was published by the Internet Engineering Task Force in February 1998 (Internet site:

<http://www.ietf.org/internet-drafts/draft-nichols-dsopdef-00.txt>) proposes a method that enables an accelerated transmission of data packets from a source node to a destination node.

For data packets to be forwarded especially fast, specific bits of what is referred to as the TOS byte in the header part of such a data packet are set. Such data packets can be divided into various classes on the basis of the bits that are set in the TOS byte. According to their class, the data packets identified with the set bits are given privileged handling in the intermediate nodes via which such data packets are transmitted from the source node to the destination node, as a result whereof, in particular, an accelerated forwarding to the next node (intermediate or destination node) is achieved.

A critical disadvantage of this method is to be seen therein that the privileged handling in the forwarding of the data packets identified with the set bits leads to considerable delays in the forwarding of the data packets to be transmitted “normally”.

Moreover, it is not only the destination addresses but also the respective TOS bytes that must be taken into consideration in every intermediate node when forwarding the data packets.

Another “Internet Draft” document authored by Y. Bernet, R. Yavatkar, P. Ford, F. Baker and L. Zhang that was published by the Internet Engineering Task Force in March 1998 (Internet site: <http://ietf.org/internet-drafts/draft-bernet-intdiff-00.txt>) presents a method that, with the assistance of the above-explained method, combines a plurality of demands made of the Internet for a guaranteed quality for the transmission of data packets into classes. A corresponding quality for the transmission of data packets is assured according to such a class.

Since a required quality for the transmission of data packets with a privileged handling thereof on the basis of the aforementioned, set bits in the TOS byte is assured, this method can be reduced to the initially explained approach. The disadvantages that were already initially mentioned therefore remain.

Substitute Page

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Another approach to offering the user of, for example, a video transmission service a bandwidth requested by the user for the transmission of data packets is disclosed by US Patent Application 5,732,078.

This discloses an arrangement of an access node to the Internet that
5 assures a bandwidth requested by the user for the transmission of data packets by re-
routing data packets onto an alternate network. The re-routing of the data packets is
thereby undertaken as follows:

When a user requests a specific bandwidth for the transmission of data
packets from his user terminal device to a destination node, the access node to which
10 the user terminal device is connected sets up a point-to-point connection to the access
node to which the destination node is connected via the alternate network offering the
requested bandwidth.

For re-routing the data packets for whose transmission the user requests an
assured bandwidth onto the alternate network, an existing routing table in the access
15 node to which the user terminal device is connected is finally modified such that, in
addition to containing the respective destination addresses of the nodes to which data
packets are respectively forwarded, it also contains the source addresses of the data
packets to be re-routed due to a quality requested by the user. According to a
modification of the appertaining, known method, a connection-individual or,
20 respectively, transmission-individual particular ("application port number") is
additionally entered into said routing table.

On the basis of the additionally stored source address and, potentially, the
aforementioned connection-individual or, respectively, transmission-individual
particular ("application port number") in the routing table, the data packets arriving at
25 such an access node can be selected according to whether they are routed over the
ordinary Internet or via the alternate network.

This method is very involved since a separate connection via the alternate network must be set up for every transmission of data packets with a requested bandwidth that is initiated by a user.

5 Additionally, the routing table must be modified in every access node to which user terminal devices are connected, namely after every transmission of data packets with a specific bandwidth initiated by a user.

Also to be considered as further outlay is that all data packets incoming at such an access node are investigated for their source address and, potentially, the
10 aforementioned connection-individual or, respectively, transmission-individual particular ("application port number"), with reference whereto a re-routing onto the alternate network can be initiated.

Finally, it is also known (Gung-Chou Lai, Ruay-Shiung Chang, "Support Qos in IP over ATM", Master Thesis [Online], June 1997 (1997-06), National Taiwan
15 University of Science and technology, Taipei, Taiwan) that a determination regarding the path via which a data packet is to be routed in a switching system can be made in the switching system on the basis of a field TOS (type of service) in the header of data packets to be switched or, respectively, to be transmitted. This measure, however, has nothing to do with a re-routing of data packets of a packet-switching network onto at
20 least one alternate network; on the contrary, the data switchings or, respectively, transmissions here respectively ensue in one and the same switching system.

The object of the invention is therefore comprised in developing a method of the species indicated in the preamble of patent claim 1 to the effect that it can be implemented with optimally little outlay and with the slightest possible effect on its
25 environment.

The object is achieved by the features recited in the characterizing part of claim 1.

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The principle of the invention is comprised in re-routing data packets of a packet-switching network for whose transmission between their source node and their destination a specific quality is requested onto at least one alternate network that
5 assures such a requested quality. This re-routing is inventively achieved in that the data packets to be routed via an alternate network are respectively identified in their source node only by a bit pattern known to the access node connected to the source node either directly or indirectly via at least one intermediate node. Upon arrival of such data packets in such an access node, the known bit patterns are respectively
10 recognized only as a result thereof, and a re-routing of the data packets identified with the known bit patterns onto an alternate network is initiated.

A critical advantage of the inventive method is to be seen therein that the table present in an access node connected to such a source node for determining the traffic paths (routing table) remains unaffected by ongoing modifications for the re-
15 routing of the data packets for whose transmission a specific quality has been requested.

It has also proven advantageous when data packets need be investigated for the known bit pattern only in the access node connected to a source node.

The invention is additionally characterized in that the traffic flow of the
20 ordinary packet-switching network is not negatively affected by the re-routing of the data packets for whose transmission a requested quality is to be assured with the assistance of the inventively simple selection according to a known bit pattern.

Further developments of the invention are recited in subclaims.

According to an advantageous development of the invention, the data
25 packets arriving in an access node directly or indirectly connected to a source node

proceeding from the source node are checked with a filter for a bit pattern known to the access node. When a known bit pattern has been recognized, then the re-routing of the data packets identified with such a bit pattern onto an alternate network is initiated. This procedure is beneficial particularly because it is no longer all data
5 packets arriving at the access node but only a significantly lower plurality of data packets, namely those coming from a source node, that must be investigated for the known bit pattern. What is also thereby avoided is that data packets that are sent from the destination node back to the source node, for example for the purpose of an acknowledge, are re-routed onto the alternate network, as a result whereof the data
10 packets are unintentionally sent in a circle.

Another advantageous development of the invention provides that the function of the filter is integrated in a table for determining the traffic paths (routing table) that is present in such an access node. This occurs with an additional entry of the known bit pattern that can produce a re-routing of a data packet identified with
15 such a bit pattern onto an alternate network. In that memory structures that are already present can be utilized in the integration of the function of the filter into the routing table, the implementation outlay of such a filter is substantially reduced.

According to a useful development of the invention, such a known bit pattern is located in the header part of a data packet to be routed via an alternate
20 network. As a result thereof, the data packets can be designationally and, thus, quickly investigated for the known bit pattern.

According to one development, the inventive method uses the respectively same bit pattern for identifying data packets to be routed via an alternate network, regardless of the respectively requested quality. The investigation of such data
25 packets for a known bit pattern is considerably simplified as a result thereof. Further, the data packets for whose transmission a respectively different transmission quality is requested are combined into a class, so that the alternate network need offer only one of the requested qualities for the transmission of these data packets.

As an alternative to the aforementioned development, the source nodes
30 can also employ different bit patterns for identifying such data packets, these different

bit patterns corresponding to the respectively requested quality. This has the advantage that such data packets can be divided into classes according to their requested quality. The alternate network can thus offer the quality corresponding to a class for the transmission of the data packets.

5 In conjunction with the above-described development, another embodiment of the invention provides that each recognized bit pattern of a data packet produces a re-routing thereof onto an alternate network with a requested quality that corresponds to the bit pattern. In other words, data packets of a class are re-routed onto one of the possible alternate networks that corresponds to the bit pattern of the
10 class and offers the quality requested by the class.

As an alternative to the aforementioned embodiment of the invention, each recognized bit pattern of a data packet can produce a re-routing thereof onto a single alternate network that assures the quality that corresponds to the bit pattern and, thus, to the class of the data packets.

15 The above-presented developments for dividing the data packets into classes corresponding to the requested quality offers [sic] the advantage that an alternate network need not explicitly offer the respectively requested quality for each transmission of data packets to be undertaken.

After recognition of such a bit pattern of a data packet to be routed via an
20 alternate network in such an access node, the re-routing of the data packet onto an alternate network can be prevented according to another advantageous development of the invention when the alternate network cannot offer the quality corresponding to the bit pattern. In this way, such an access node can defend such an alternate network against data packets that lead to an overload.

25 An advantageous development of the invention provides that the source node that intends to send data packets communicates a message to its destination node via the packet-switching network with respect to the data packets to be routed via an alternate network and potentially waits for an acknowledge from the destination node. A secured connection setup of a point-to-point connection between the source and the
30 destination node is achieved as a result thereof.

Another development of the invention is comprised therein that the access node directly or indirectly connected to such a source node sends a message with respect to the quality to be assured as requested by the source node to the network node of an alternate network and potentially waits for an acknowledge therefrom.

5 This procedure serves the purpose of a dependable offering of the requested quality by the network node of the alternate network.

A further advantageous development of the inventive method can be seen therein that the method is applied to a network constellation wherein such an alternate network is formed in that at least one logical channel of the packet-switching network
10 is reserved for the data packets to be transmitted with an assured quality. I.e., an alternate network is composed of one or more logical channels of the packet-switching network. This represents an especially cost-beneficial solution since additional connecting lines for the alternate network are thereby eliminated.

An exemplary embodiment of the invention is explained in greater detail
15 below with reference to a drawing.

The Figure shows a portion of a communication network on which the inventive method can be applied by way of example.

Accordingly, the Figure shows two source nodes U1, U2 of a plurality of conceivable source nodes and one of many possible destination nodes Z. The two
20 source nodes U1, U2 are connected to an access node ZK1 and the destination node Z is connected to an access node ZK2, either directly or indirectly via respectively one or more intermediate nodes ZW1, ZW2, as indicated in the Figure by the intermediate node ZW1 shown with broken lines between U1 and ZK1 and the intermediate node ZW2 shown with broken lines between ZK2 and Z. The access nodes ZK1 and ZK2
25 belong to a packet-switching network PN and respectively form an access to one or more alternate networks AN1 through ANn at the same time. Such a packet-switching network could, for example, be the Internet, within which mainly variable-length data packets are transmitted, or could also be an ATM network (asynchronous transfer mode), within which fixed-length data packets (ATM cells) are usually
30 transported. Given the pre-condition that it assures a requested quality for the data

transmission, such an alternate network can be formed by an arbitrary network type, for example a line-switching network, a packet-switching network or an ATM network. Such an alternate network can also be composed of at least one logical channel of the packet-switching network PN.

5 Normally, data packets are sent to an access node, for example ZK1, to the packet-switching network PN from the source node either directly (for example, proceeding from U2) or indirectly (for example, proceeding from U1) via one or more intermediate nodes, for example ZW1. From the access node, they are sent via the packet-switching network to another access node, for example ZK2, and are sent from
10 the latter to a destination node Z either directly or via one or more intermediate nodes, for example ZW2. A transmission of data packets in the opposite direction is likewise conceivable, i.e. with the destination node as source node and with the source node as destination node.

 In view of the overall communication network, the indicated source node
15 and the destination node can be respectively interpreted as intermediate node of end node of the communication network. Such end nodes can thereby be viewed as computer of a service vendor or as a data terminal device of a user.

 In the present example, a user would like to have his data terminal device, for example U1, receive an Internet service, for example video on demand, from a
20 computer, for example Z, of a service vendor. For the transmission of data packets of such a service, a certain quality in the form of a specific bandwidth must be assured at the network side. The bandwidth needed for the service can be offered by re-routing the data packets belonging to such a service onto an alternate network that assures the requested quality.

25 Accordingly, the source node, for example U1, sends a message to the access node ZK1 either directly or indirectly via at least one intermediate node, for example ZW1, this message containing a request in the form re-routing data packets for whose transmission a specific quality, for example bandwidth, is required. In the form of a message, the access node communicates the bit pattern with which the data
30 packets to be routed onto the alternate network are identified to the source node U1

and additionally acknowledges the message with the request sent from the source node U1. Further, the access node sends a message about the quality requested by the source node to the network node (not shown in the Figure) of an alternate network, for example AN1, and potentially waits for the acknowledge thereof. The source node
5 U1 identifies the data packets that are to be transmitted via the alternate network with an assured quality with the bit pattern obtained from the access node.

As an alternative thereto, such a bit pattern can be known to the source node U1 according to a corresponding implementation without having to be previously informed thereof by the access node. Such a bit pattern is normally located
10 in the header of such a data packet. In conjunction with the Internet, the bit pattern can be found in what is referred to as the TOS byte.

Before the source node U1 now sends the data packets -- for whose transmission a requested quality is to be assured and which are therefore to be routed via an alternate network -- in the direction of the access node, it is possible that the
15 source node communicates a message with respect to the [...] to be routed via the alternate network to its destination node data packets [sic] via the ordinary packet-switching network PN and potentially waits for an acknowledge from the destination node.

In the access node, the data packets arriving from the source node U1 are
20 checked for the bit pattern with a filter (not shown in the Figure). The data packets wherein the bit pattern was recognized are re-routed onto the alternate network. Alternatively thereto, the function of the filter can be integrated into the routing table to be conceived of as being present in the access node in that the bit patterns that produce a re-routing of a data packet identified with the bit pattern onto the alternate
25 network are additionally entered into the routing table.

Parallel to the above-described scenario with respect to the source node U1, the same scenario can be initiated by one or more other source nodes, for example U2.

Regardless of the respectively requested quality, the source node U2
30 possibly employs the same bit pattern for identifying the data packets that are to be

sent via an alternate network, for example ANn. Alternatively thereto, the source node U2 can use a bit pattern that corresponds to the requested quality for identifying the data packets to be routed via the alternate network.

5 In this way, the data packets sent in the direction of the access node from, under certain circumstances, a plurality of source nodes can be combined into classes according to the quality required for their transmission.

10 When the data packets to be re-routed in the access node can be re-routed onto a plurality of alternate networks, one bit pattern is defined [...] each alternate network. In this way, the data packets can be re-routed onto an alternate network corresponding to the bit pattern that assures the requested quality.

When only one alternate network is available for the re-routing of the data packets identified with a bit pattern, the quality to be assured by the alternate network can be offered according to the bit pattern of the identified data packets.

15 The access node can deny re-routing onto the alternate network for data packets that are identified with a bit pattern that does not correspond to the quality offered by the alternate network.

Mutually independent actions of the above-described scenarios can also run in an arbitrary sequence.

Patent Claims

1. Method for re-routing data packets of a packet-switching network (PN) onto at least one alternate network (AN1, ..., ANn) that assures a quality requested for these data packets, whereby the packet-switching network and the at least one
5 alternate network form sub-networks of a communication network that is composed of at least one source node (U1, U2) and at least one destination node (Z) that are respectively connected to an access node (ZK1, ZK2) either directly or indirectly via at least one intermediate node, said access node being capable of setting up a
10 connection both to the packet-switching network as well as to an alternate network, and between which (U1, U2; Z) data packets can be transmitted, characterized in that, for the re-routing of the data packets via an alternate network (AN1, ..., ANn) assuring the quality requested for these data packets, the appertaining data packets are identified in their source node only by a respective bit pattern known from such a re-
15 routing to the access node (ZK1) that is connected to the source node either directly or indirectly via at least one intermediate node (ZW1); and in that, upon arrival of such data packets respectively identified by the appertaining bit pattern in the access node, a re-routing of the data packets identified with the known bit pattern onto an alternate network (AN1, ..., ANn) is initiated based only on the recognition of the appertaining
20 bit pattern.

2. Method according to claim 1, characterized in that a filter in such an access node (ZK1, ZK2) checks the data packets arriving from a source node (U1, U2) for a bit pattern known to the access node (ZK1, ZK2) and, when a known bit pattern was recognized, initiates the re-routing of the data packets identified with this bit
25 pattern onto an alternate network (AN1, ..., ANn).

3. Method according to claim 1, characterized in that the access node (ZK1, ZK2) connected to such a source node (U1, U2) either directly or indirectly via at least one intermediate node (ZW1, ZW2) contains a table for determining the traffic
5 paths into which the function of the filter is integrated, in that it additionally contains the bit patterns that can produce a re-routing of a data packet identified with such a bit pattern onto an alternate network (AN1, ..., ANn).

4. Method according to one of the preceding claims, characterized in that such a bit pattern is located in a header of a data packet to be routed via an alternate
10 network (AN1, ..., ANn).

5. Method according to one of the preceding claims, characterized in that the same bit pattern is employed by all source nodes (U1, U2) regardless of the respectively requested quality.

6. Method according to one of the claims 1 through 4, characterized in that
15 the source nodes (U1, U2) employ bit patterns corresponding to the respectively requested quality.

7. Method according to claim 6, characterized in that each recognized bit pattern of a data packet produces a re-routing thereof onto an alternate network (AN1, ..., ANn) corresponding to the bit pattern with a specific quality.
20

8. Method according to claim 6, characterized in that each recognized bit pattern of a data packet produces a re-routing thereof onto an alternate network (AN1, ..., ANn) with a quality corresponding to the bit pattern.

9. Method according to one of the preceding claims, characterized in that, after recognition of such a bit pattern of a data packet to be routed via an alternate
25 network (AN1, ..., ANn) in such an access node (ZK1, ZK2), the re-routing of the data packet onto an alternate network (AN1, ..., ANn) can be prevented when the alternate network (AN1, ..., ANn) cannot offer the quality corresponding to the bit pattern.

10. Method according to one of the preceding claims, characterized in that the source node that intends to send data packets communicates a message via the packet-switching network (PN) to its destination node (Z) with respect to the data
5 packets to be routed via an alternate network (AN1, ..., ANn) and potentially waits for an acknowledge from the destination node (Z).

11. Method according to one of the preceding claims, characterized in that the access node (ZK1, ZK2) connected to such a source node (U1, U2) either directly or indirectly via at least one intermediate node (ZW1, ZW2) sends a message with
10 respect to the assured quality requested by the source node (U1, U2) to the network node of such an alternate network (AN1, ..., ANn) and potentially waits for an acknowledge thereof.

12. Method according to one of the preceding claims, characterized in that it is applied to a network constellation wherein such an alternate network (AN1, ...,
15 ANn) is formed in that at least one logical channel of the packet-switching network (PN) is reserved for the data packets to be transmitted with an assured quality.

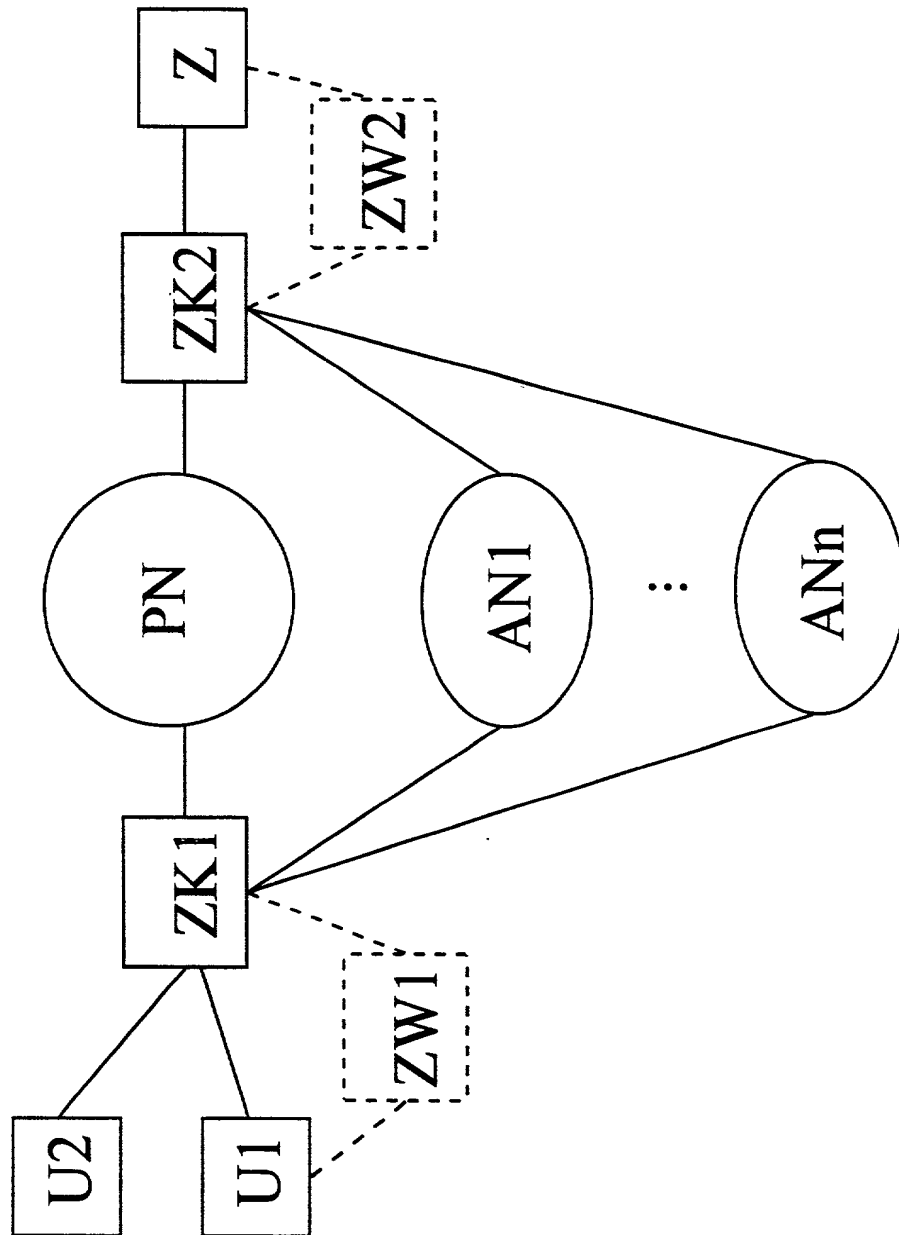
Abstract

**METHOD FOR RE-ROUTING DATA PACKETS OF A PACKET SWITCHING
NETWORK ONTO AT LEAST ONE ALTERNATIVE NETWORK**

Data packets of a packet-switching network (PN) for whose transmission a
5 requested quality is to be assured are routed via at least one alternate network (AN 1,
..., ANn). In their source node (U1, U2), such data packets are respectively identified
by a bit pattern known [...] access node (ZK1) connected to the source node either
directly or indirectly via at least one intermediate node (ZW1), said bit pattern being
respectively recognized upon arrival of such data packets in the access node, as a
10 result whereof a re-routing of the data packets identified with the known bit pattern
onto an alternate network is initiated.

Figure

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Verfahren Zur Umleitung Von Datenpaketen Auf Ein Alternatives Netz

deren Beschreibung

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Prior foreign applications
Priorität beansprucht

Priority Claimed

198 34 977.7 Germany 03 August 1998
(Number) (Country) (Day Month Year Filed)
(Nummer) (Land) (Tag Monat Jahr eingereicht)

☒ ☐
Yes No
Ja Nein

(Number) (Country) (Day Month Year Filed)
(Nummer) (Land) (Tag Monat Jahr eingereicht)

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(Anmeldeseriennummer)

(Filing Date)
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aufgegeben)

(Status)
(patented, pending,
abandoned)

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(Anmeldeseriennummer)

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